

## Strong Synchronized System

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**Abstract.** In this paper introduced the notion of a strong synchronized system; that is a synchronized system whose there is unique finite path in Fischer cover labeled synchronizing block. We aim to introduce a class of synchronized systems containing sofics. Every irreducible sofic shift is an strong synchronized.

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### 1 Introduction

One of the most studied dynamical systems is a subshift of finite type (SFT). An SFT is a system whose set of forbidden blocks is finite ([6]). Equivalently, an SFT  $X$  is a subshift whose any block of length greater than a certain number  $M$  is synchronizing; that is, if  $m$  is *any* block with  $|m| \geq M$  and if  $v_1m$  and  $mv_2$  are both blocks of  $X$ , then  $v_1mv_2$  is a block of  $X$ . If an irreducible system has at least one synchronizing block, then it is called a *synchronized system* and examples are *sofics*: factors of SFT's.

For a synchronized system, Fiebig in ([2]) prove that there is some finite path  $e$  in Fischer cover labeled  $m$  terminating in  $\alpha$  such that

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$m \in F_-(\alpha)$ . In the other hand,  $\text{cardinal}\{e \in \mathcal{E}_{X_0^+} : t(e) = \alpha\} \geq 1$ . In this note we consider shift space that  $\text{cardinal}\{e \in \mathcal{E}_{X_0^+} : t(e) = \alpha\} = 1$ .

## 2 Background and Definitions

This section is devoted to the very basic definitions in symbolic dynamics. The notations have been taken from ([6]) and the proof of the relevant claims in this section can be found there. Let  $\mathcal{A}$  be an alphabet, that is a non-empty finite set. The full shift  $\mathcal{A}$ -shift denoted by  $\mathcal{A}^{\mathbb{Z}}$ , is the collection of all bi-infinite sequences of symbols in  $\mathcal{A}$ . A block over  $\mathcal{A}$  is a finite sequences of symbols from  $\mathcal{A}$ . It is convenient to include the sequence of no symbols, called the *empty block* and denoted by  $\varepsilon$ . If  $x$  is a point in  $\mathcal{A}^{\mathbb{Z}}$  and  $i \leq j$ , then we will denote a block of length  $j - i$  by  $x_{[i,j]} = x_i x_{i+1} \dots x_j$ . If  $n \geq 1$ , then  $u^n$  denotes the concatenation of  $n$  copies of  $u$ , and put  $u^0 = \varepsilon$ . Let  $\mathcal{F}$  be the collection of all forbidden blocks over  $\mathcal{A}$ . The complement of  $\mathcal{F}$  is the set of admissible blocks or just blocks in  $X$ . For any such  $\mathcal{A}^{\mathbb{Z}}$ , define  $X_{\mathcal{F}}$  to be the subset of sequences in  $\mathcal{A}^{\mathbb{Z}}$  not containing any block in  $\mathcal{F}$ . A shift space is a closed subset  $X$  of a full shift  $\mathcal{A}^{\mathbb{Z}}$  such that  $X = X_{\mathcal{F}}$ . For some collection  $\mathcal{F}$  of forbidden blocks over  $\mathcal{A}$ .

Let  $W_n(X)$  denote the set of all admissible  $n$  blocks. The language of  $X$  is the collection  $W(X) = \cup_n W_n(X)$ . A shift space  $X$  is *irreducible* if for every ordered pair of blocks  $u, v \in W(X)$  there is a block  $w \in W(X)$  so that  $uwv \in W(X)$ . It is *mixing* if for every ordered pair  $u, v \in W(X)$ , there is an  $N$  such that for each  $n \geq N$  there is a block  $w \in W_n(X)$  such that  $uwv \in W(X)$ . A shift space  $X$  is called a *shift of finite type* SFT if there is a finite set  $\mathcal{F}$  of forbidden blocks such that  $X = X_{\mathcal{F}}$ . An *edge shift* denoted by  $X(G)$ , is a shift space that consist of all bi-infinite walks in a directed graph  $G$ . Each edge  $e$  initiates at a vertex denoted by  $i(e)$  and terminates at a vertex  $t(e)$ .

A *labeled graph*  $\mathcal{G}$  is a pair  $(G, \mathcal{L})$  where  $G$  is a graph with edge set  $\mathcal{E}$ , and the labeling  $\mathcal{L} : \mathcal{E} \rightarrow \mathcal{A}$ . A *sofic shift*  $X_{\mathcal{G}}$  is the set of sequences obtained by reading the label of walks on  $G$ .

$$X_{\mathcal{G}} = \{\mathcal{L}_{\infty}(\varepsilon) : \varepsilon \in X_{\mathcal{G}}\} = \mathcal{L}_{\infty}(X_{\mathcal{G}}).$$

We say  $\mathcal{G}$  is a presentation of  $X_{\mathcal{G}}$ . The follower set of a vertex  $m$  of  $G$

is  $F_-(m) = \{\mathcal{L}\text{-label of all finite paths terminating at } m\}$ . Every SFT is *sofic* ([6, Theorem 3.1.5]), but the converse is not true. A labeled graph  $\mathcal{G} = (G, \mathcal{L})$  is *right-resolving* if for each vertex  $I$  of  $G$  the edges starting at  $I$  carry different labels.

Let  $X$  be a shift space and  $w \in W(X)$ . The follower set  $F_+(w)$  of  $w$  is defined by  $F_+(w) = \{v \in W(X) : wv \in W(X)\}$  (resp.  $F_-(w)$ ). Let  $x \in W(X)$ . Then,  $x_+ = (x_i)_{i \in \mathbb{Z}^+}$  (resp.  $x_- = (x_i)_{i < 0}$ ) is called *right* (resp. *left*) infinite X-ray. For a left infinite X-ray, say  $x_-$  its follower set is  $w_+(x_-) = \{x^+ \in X^+ : x_-x^+ \in X\}$ . Consider the collection of all follower sets  $w_+(x_-)$  as the set of vertices of a graph  $X^+$ . There is an edge from  $I_1$  to  $I_2$  labeled  $a$  if and only if there is an X-ray  $x_-$  such that  $x_-a$  is an X-ray and  $I_1 = w_+(x_-), I_2 = w_+(x_-a)$ . This labeled graph is called the *Krieger graph* for  $X$ . A block  $m \in W(X)$  is *synchronizing* if whenever  $um$  and  $mv$  are in  $W(X)$ , we have  $umv \in W(X)$ . An irreducible shift space  $X$  is a synchronized system if it has synchronizing block. A block  $m \in W(X)$  is *half synchronizing* if there is a left transitive point  $x \in X$  such that  $x_{[-|m|+1, 0]} = m$  and  $w_+(x_{(-\infty, 0]}) = w_+(m)$ , That we denote by

$$(x, m). \tag{1}$$

If  $X$  is a half synchronized system with half synchronizing  $m$ , the irreducible component of the Krieger graph containing the vertex  $w_+(m)$  is denoted by  $X_0^+$  and is called the *right Fischer cover* of  $X$ .

Let  $X$  be a shift space. The *entropy* of  $X$  is defined by

$$h(X) = \lim_{n \rightarrow \infty} \frac{1}{n} \log |W_n(X)|.$$

For any synchronized system  $X$ , we define the *synchronized entropy*  $h_{\text{syn}}$  to be

$$h_{\text{syn}}(X) = \limsup_n \frac{1}{n} \log(\text{cardinal}\{a \in W_n(Y) : mam \in W(X)\}),$$

where  $m$  is an arbitrary synchronizing block in  $W(X)$ . A shift space  $X$  is *almost sofic* if there are sofic shifts  $X_n \subseteq X$  such that  $\lim_{n \rightarrow \infty} h(X_n) = h(X)$ .

### 3 Strong Synchronized Systems

We aim to introduce a class of synchronized systems containing sofics. Let  $X$  be a shift space and  $R(X) = \overline{\text{Per}X}$  and set  $S(X)$  to denote the set of synchronizing blocks for  $R(X)$ . For  $s, t \in S(X)$  we write  $s \sim t$  whenever there are blocks  $u, v \in W(R(X))$  such that  $sut, tvs \in W(R(X))$ . Then,  $\sim$  is an equivalence relation in  $S(X)$ . Note that  $s \sim t$  if only if there is an  $x \in R(X) = \overline{\text{Per}X}$  such that  $s, t \subseteq x$ . Let  $x \in R(X)$  and for integers  $p$  and  $s$  with  $p \leq s$ , set

$$\text{gap}(x_{[p,q]}, x_{[s,t]}) = \begin{cases} 0 & s \leq q \\ s - q & \text{otherwise} \end{cases}$$

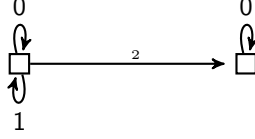
and call it the *gap* between two blocks  $x_{[p,q]}$  and  $x_{[s,t]}$ . Pick  $\alpha \in S(X)/\sim$ ,  $p < s$ ,  $q \leq t$  and  $\{u = x_{[p,q]}, v = x_{[s,t]}\} \subseteq \alpha$ . If the only minimal synchronizing blocks in  $x_{[p,t]}$  are  $u$  and  $v$ , then call  $u$  and  $v$  the *consecutive minimal pairs of  $\alpha$  in  $x_{[p,t]} \subset x$* . In ([2])  $X_{(\alpha,0)}$  was defined to be the set of elements of  $x \in R(X)$  satisfying the following.

- i. If  $i \in \mathbb{Z}$ , then there are  $m, m' \in \alpha$  such that  $m \subseteq x_{(-\infty, i]}$ ,  $m' \subseteq x_{[i, +\infty)}$ .
- ii.  $|\{u \in \alpha : u \subseteq x\}| < \infty$ .
- iii. There is  $M > 0$  such that if  $u$  and  $v$  are the consecutive minimal pairs of  $\alpha$  in  $x$ , then  $\text{gap}(u, v) \leq M$ .

**Example 3.1.** Let  $G$  be the graph in Figure 1 and  $X = X(G)$ . Then,  $X_{(\alpha,0)} = \{0, 1\}^{\mathbb{Z}}$ .

We associate to  $X_{(\alpha,0)}$  an irreducible graph  $\Gamma_\alpha$ . To do so for each  $m \in \alpha$ , let  $F(m) = \{u \in W(X_{(\alpha,0)}) : mu \in W(X_{(\alpha,0)})\}$  and let the vertices of  $\Gamma_\alpha$  to be  $\{F(m) : m \in \alpha\}$ . Assign an edge labeled  $a$  from  $F(m)$  to  $F(m')$  when  $a \in F(m)$  and  $F(ma) = F(m')$ . Then,  $\Gamma_\alpha$  is a minimal right resolving cover and is called *Fischer cover* of  $\overline{X_{(\alpha,0)}}$ .

**Definition 3.2.** A block  $m \in W(X)$  is called strong synchronizing if there is an element  $\alpha \in S(X)/\sim$  such that whenever there are two finite paths  $e, e'$  in  $\Gamma_\alpha$  labeled  $m$ , then  $e = e'$ .



**Figure 1:** The graph  $G$ .

Call an irreducible shift space  $X$  *strong synchronized* if it has a strong synchronizing block. In the sequel, we will show that

$$\text{sofics} \subsetneq \text{strong synchronized systems} \subsetneq \text{synchronized systems}. \quad (2)$$

Clearly, any strong synchronizing block is a synchronizing block. The next example shows that the second inclusion in (2) is not equality.

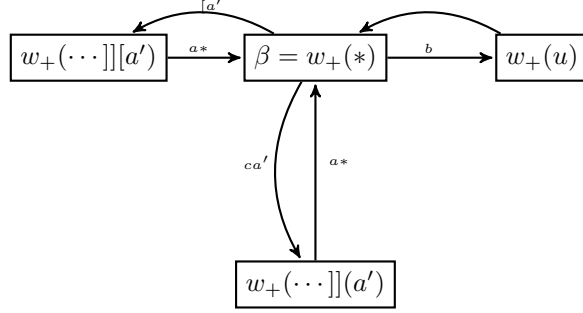
**Example 3.3.** Let  $X$  be the Dyke system. Add a new symbol  $*$  (which will be a synchronizing block) to the set of four brackets and let  $S$  be the subshift which consists of all bi-infinite sequences of these five symbols such that any finite subblock which does not contain a  $*$  obeys the standard brackets rules ([2]). We claim that  $S$  is not strong synchronized system.

Let  $u$  be a strong synchronizing block of  $S$ . So that is a synchronizing block and  $* \subseteq u$ . We can write  $u = a * b$  where  $a, b \in W(S)$ . Since  $*$  is a synchronizing block, there is a unique vertex  $\beta$  of Fischer cover  $S_0^+$  with this properties  $* \in F_-(\beta)$  and  $w_+(\beta) = w_+(*)$  ([2, Page 146]). Since  $*b \in W(S)$ , so there is a cycle  $C$  containing  $b$  and passing through  $w_+(*)$  and  $w_+(u)$ . See Figure 2.

Pick  $a' \in W(S)$  such that  $[a', (a' \in W(S) \text{ and for each } x_- \in S^-, x_- a' a \in S^-)$ . Set  $x_- := \cdots ]][a', y_- := \cdots ])(a'$ . Since  $w_+(x_-) \neq w_+(y_-)$  and

$$w_+(x_- a *) = w_+(y_- a *), w_+(*[a') = w_+(x_-), w_+(*a') = w_+(y_-),$$

so there are finite pathes  $e, e'$  labeled  $a * b = u$  and terminating at  $w_+(u)$ . This show that synchronized system  $S$  is not strong synchronized.



**Figure 2:** A subgraph of Fischer cover  $S$ .

If  $m$  is a synchronizing block of  $X$ , then there is a unique vertex  $I$  of  $\Gamma_\alpha$  with the property  $m \in F_-(I)$  as in  $w_+(u)$  in Figure 2. But  $m$  is a strong synchronizing block if and only if there are two unique vertexes  $I$  and  $J$  of  $\Gamma_\alpha$  such that  $m \in F_-(I)$  and  $m \in F_+(J)$ . Note that if  $X$  is a synchronized system, then  $\{\alpha\} = S(X)/\sim$  and  $\Gamma_\alpha = X_0^+$ .

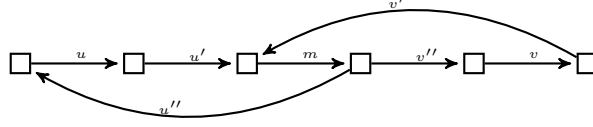
**Proposition 3.4.** *Every irreducible sofic shift is an strong synchronized.*

**Proof.** Let  $m$  be a synchronizing block of  $X$ . There is a finite path  $\pi$  in Fischer cover  $X_0^+$  labeled  $m$ . Set  $i(\pi) := I$  and  $t(\pi) := J$ . Since labeling of finite graph  $X_0^+$  is right resolving, so the labeling of  $X_0^+$  is left closing ([2, Theorem 3.2]) and so  $m$  is a strong synchronizing block of  $X$ .  $\square$  The next example shows that the converse of the above proposition is not necessarily true and so (2) is sort out completely.

**Example 3.5.** Set  $1_\beta := 21010^210^31\dots$ . Since  $1_\beta$  is not eventually periodic so by ([3, Theorem 2.4.2])  $\beta$ -shift is not sofic. Since 2 is a strong synchronizing block for  $\beta$ -shift, so it is a strong synchronized.

Set  $S_t(X)$  to denote the set of strong synchronizing blocks for  $X$ . If  $m \in S_t(X)$  and  $\alpha = [m]$ , then for each  $a \in W(X_{(\alpha,0)})$ , let  $C'_a$  be a cycle in  $\Gamma_\alpha$  containing  $a$  and passing through  $w_+(m)$ . Let  $u, v \in W(X_{(\alpha,0)})$ . Set the *distance* between two blocks  $u$  and  $v$ , to be 0 when  $u = v$  and

$$\min\left\{\frac{1}{2}(|u'mu''|+|v'mv''|) : \mathcal{L}(C'_{uu'mu''}) = uu'mu'', \mathcal{L}(C'_{vv'mv''}) = vv'mv''\right\}$$



**Figure 3:** The distance of  $u$  and  $v$ .

otherwise. Denote this distance by  $d_m(u, v)$  see Figure 3. It is not hard to see that  $(W(X_{(\alpha,0)}), d_m)$  is a metric space. Note that the definition of  $d_m$  is dependent on the particular choice of  $m$ . Let  $m_1, m_2 \in \alpha \cap S_t(X)$ . Then, for each  $u \in W(X_{(\alpha,0)})$   $d_{m_1}(m_2, u) = d_{m_1}(m_1, u) + d_{m_2}(m_1, m_2) + |m'_1| + |m_2| - |m'_2|$ .

**Lemma 3.6.** *If  $uv \in W(X_{(\alpha,0)})$  and  $m \in S_t(X)$ , then*

- i.  $d(v, m) \leq \frac{1}{2}|u| + d(uv, m)$ .
- ii.  $d(u, m) \leq \frac{1}{2}|v| + d(uv, m)$ .

**Proof.** Let  $d(v, m) = \frac{1}{2}(|v'| + |v''| + 2|m| + 2|m'|)$  and

$$d(uv, m) = \frac{1}{2}(|a| + |b| + 2|m| + 2|m'|) \quad (3)$$

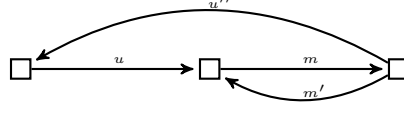
for some  $v', v'', a, b \in W(X_{([m],0)})$ . By definition,  $|a| + |m| + |b| + |u| \geq |v'| + |m| + |v''|$  So  $d(v, m) \leq \frac{1}{2}(|a| + |b| + |u| + 2|m| + 2|m'|)$  and by (3),  $d(v, m) \leq \frac{1}{2}|u| + d(uv, m)$ . Similar reasoning works for (ii).  $\square$

**Lemma 3.7.** *labelaub*

- i. Let  $m \in S_t(X)$  and  $d(u, m) = \frac{1}{2}(|u'| + |u''| + 2|m| + 2|m'|)$ .

(a) *If  $u' = a'a''$ ,  $u'' = b'b''$ , then,*

$$d(b''ua', m) = \frac{1}{2}(|b'| + |a''| + 2|m| + 2|m'|) \leq d(u, m).$$



**Figure 4:** Distance of  $u$  and  $m$  when  $um \in W(X_{(\alpha,0)})$ .

(b) If  $u' = \varepsilon$  and  $u'' \neq \varepsilon$ , then

$$d(u''umm', m) = |m| + |m'| = d(u, m) - \frac{1}{2}|u''|.$$

(c) If  $u' \neq \varepsilon$  and  $u'' = \varepsilon$ , then

$$d(m'muu', m) = |m| + |m'| = d(u, m) - \frac{1}{2}|u'|.$$

(d) If  $u' = u'' = \varepsilon$ , then  $d(m'mumm') = d(u, m) = |m| + |m'|$

ii. There are nonempty blocks  $a$  and  $b$  such that  $d(aub, m) \leq d(u, m)$ .

**Proof.** (a) Let  $C$  be a cycle containing  $b''ua'$  and  $m$ . It suffice to show that

$$|C| \geq |b''| + |u| + |a'| + |a''| + |m| + |b'|.$$

Let  $\mathcal{L}(C) = b''ua'w'mw''$ . Then, by definition,

$$|a'| + |w'| + |m| + |w''| + |b''| \geq |u'| + |m| + |u''| = |a'| + |a''| + |m| + |b'| + |b''|$$

and we are done.

(b) Let  $\pi_a$  be a path in  $\Gamma_\alpha$  and labeled  $a$ . There is a finite path  $\pi_{u''umm'}$  terminating at  $i(\pi_m)$  and starting at  $t(\pi_m)$ , so (b) is trivial. See Figure 4.

Similar reasoning works for (c) and (d).

The part of (ii) follows from (i).  $\square$

**Proposition 3.8.** Let  $m \in S_t(X)$  and  $r > 0$ . Then,

$$N_{m,r} := N_r(m) \cup \{u \subseteq v : v \in N_r(m)\}$$

is a languages of a shift space.



**Proof.** We prove the proposition by showing that if  $u \in N_{m,r}$ , then

- i. every subblock of  $u$  belongs to  $N_{m,r}$ .
- ii. there are nonempty blocks  $a$  and  $b$  in  $N_{m,r}$  so that  $aub \in N_{m,r}$ .

Definition of  $N_{m,r}$  implies that (i) is trivial. To prove (ii), let  $u \in N_r(m)$ . By Lemma ??, there are nonempty blocks  $a, b$  such that  $d(aub, m) \leq d(u, m)$ . So  $aub \in N_r(m)$  and we are done.  $\square$  Let  $X_{(m,r)}$  be a shift space such that  $N_{m,r}$  is its language. We can associate to  $X_{(m,r)}$  an irreducible graph  $\Gamma_m$ . If  $u \in N_r(m)$ , then  $d(u, m) = \frac{1}{2}(|u'| + |u''| + 2|m| + 2|m'|)$  for some  $u', u'' \in W(X_{(\alpha,0)})$ .

Let  $C_u$  be the cycle in  $\Gamma_\alpha$  labeled  $uu'mu''$  and passing through  $w_+(m)$ . Then,  $\Gamma_m$  consist of all  $C_u$  when  $d(u, m) < r$ . It is easy to see that  $W(X_{(m,r)}) \subseteq W(\mathcal{L}(\Gamma_m))$ . Now suppose  $\pi_u$  be a finite path in  $\Gamma_m$  labeled  $u$ . So there is a cycle  $C_v$  in  $\Gamma_m$  labeled  $vv'mv''$  such that  $u \subseteq vv'mv''$ . Since  $d(v''vv', m) = |m| + |m'| \leq d(v, m) < r$ , thus  $v''vv' \in N_r(m)$  and so  $u \in N_{m,r}$  or  $W(\mathcal{L}(\Gamma_m)) \subseteq W(X_{(m,r)})$ . Thus  $X_{(m,r)} = \mathcal{L}(\Gamma_m)$ . Also by definition of  $\Gamma_m$ , it is a right resolving graph and follower separated. Which set over claim by ([9, page 3563]) and so  $\Gamma_m$  is Fischer cover of  $X_{(m,r)}$ .

**Proposition 3.9.** *Let  $X$  be a shift space and  $m \in S_t(X)$ . Then,  $X_{(m,r)}$  is a sofic.*

**Proof.** Set  $\alpha := [m]$  and  $Z := X_{(\alpha,0)}$ . Let

$$\{(i, j) \in \mathbb{Z}^2 : i, j \geq 0 \text{ and } i+j < r-2|m|+2|m'|\} = \{(i_1, j_1), \dots, (i_k, j_k)\}$$

and for each  $1 \leq l \leq k$  let  $|\{u' \in W(Z) : u'm \in W(Z) \text{ and } |u'| = i_l\}| = n'_l$ ,  $|\{u'' \in W(Z) : mu'' \in W(Z) \text{ and } |u''| = j_l\}| = n''_l$ . Then, by the fact that if  $u \in N_r(m)$  then  $|u'| + |u''| < r - 2|m| - 2|m'|$  for some  $u', u'' \in W(X_{(\alpha,0)})$ , so

$$|\{C_u : u \in N_r(m)\}| \leq n'_1 n''_1 + \dots + n'_k n''_k.$$

Thus  $\Gamma_m$  s a finite graph or  $X_{(m,r)}$  is a sofic.  $\square$

**Remark 3.10.** Note that as in half synchronized case in ([2]), we can define  $X_{(m,r)}$ . For this let  $X$  be a shift space. For  $m, m' \in S_t(X)$  we write  $m \sim' m'$  when there are  $u \in W(R(X))$  such that  $m, m' \subseteq \mathcal{L}(C_u)$ . Then,  $\sim'$  is an equivalence relation in  $S_t(X)$ .

**Corollary 3.11.** *i. Every sofic shift is a bounded metric space.*

*ii. Let  $X$  be an irreducible bounded metric space and  $m \in S_t(X)$ . Then, there is  $r > 0$  such that  $X = X_{(\alpha, 0)} = X_{(m, r)}$ .*

**Proof.** (i) Let  $X$  be a sofic and  $m \in S_t(X)$ . There is a cycle  $C$  in  $X_0^+$  such that for any  $I \in \mathcal{V}_{X_0^+}$ ,  $C$  passing through  $I$ . Let length of  $C$  be  $M$  and  $u \in W(X) = W(X_{(\alpha, 0)})$  where  $\alpha = [m]$ . Then,  $|u'|, |u''|, |m'| \leq M$ . So  $d(u, m) < 4|M| + |m|$  or  $u \in N_r(m)$  where  $r := 4|M| + |m|$ . So  $X \subseteq N_r(m)$ .

(ii) There is  $r > 0$  such that  $X \subseteq N_r(m)$ . Thus  $W(X) \subseteq W(X_{(m, r)}) \subseteq W(X_{(\alpha, 0)})$  and since  $W(X_{(\alpha, 0)}) \subseteq W(X)$ , so  $X = X_{(\alpha, 0)} = X_{(m, r)}$ .  $\square$   
Note that  $X_{(\alpha, 0)}$  is clearly a subshift of  $R(X)$ , but generally not closed. see Example 3.3.

**Corollary 3.12.** *Let  $m \in S_t(X)$  and every cycle in  $\Gamma_m$  containing  $m$ . Then,  $X$  is a sofic if and only if  $X$  is an SFT.*

**Proof.** Let  $X$  be a sofic system. So there is  $r > 0$  such that  $X = X_{(m, r)}$ . Similar to Proposition 3.9, there is  $k \in \mathbb{N}$  such that

$$\{C_u : u \in W(X_{(m, r)})\} := \{C_{u_1}, C_{u_2}, \dots, C_{u_k}\}.$$

Set  $M := \max\{|C_{u_1}|, |C_{u_2}|, \dots, |C_{u_k}|\}$  and let  $u \in W(X_{(m, r)})$  where  $|u| \geq 2M$ . By definition of  $\Gamma_m$ , there must be at least one  $m \subseteq u$  and such  $u$  is essentially synchronizing. As a result, any block of length  $2M$  in  $W(X_{(m, r)})$  is a synchronizing block and so  $X = X_{(m, r)}$  is an SFT ([6, Theorem 2.1.8]).  $\square$  The conclusion of Corollary 3.12 is not true when there is a cycle in  $\Gamma_m$  labeled  $m$ .

Thomsen in ([9]) considers a synchronized component  $X$  of a general subshift and proves that

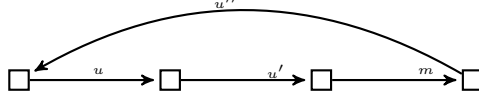
$$\sup\{h(A) : A \subseteq X \text{ is an irreducible SFT}\} = h(X_0^+) = h_{\text{syn}}(X). \quad (4)$$

This was extended to half synchronized component in ([6]) by showing that

$$\sup\{h(A) : A \subseteq X \text{ is an irreducible sofic}\} = h(X_0^+) = h_{\text{hsyn}}(X).$$

Now we will introduce this notion to strong synchronized.

Recall that if  $m \in S_t(X)$ ,  $\alpha := [m]$  and  $u \in W(X)$ , then  $C_u$  is a cycle in  $\Gamma_\alpha$  labeled  $uu'mu''$  and passing through  $w_+(m)$ .



**Figure 5:** The subgraph of  $\Gamma_m$ .

**Proposition 3.13.** *Let  $m \in S_t(X)$ . Set  $\alpha := [m]$ ,  $Z := X_{(\alpha, 0)}$ . Set*

$$\{C_u : u \in W(Z)\} := \{C_{u_1}, C_{u_2}, \dots\}$$

and

$$A_{(m, 1)} := X(C_{u_1}), A_{(m, 2)} := X(C_{u_1} \cup C_{u_2}), \dots$$

Then,  $\overline{X_{(\alpha, 0)}} = \overline{\cup_{n \in \mathbb{N}} A_{(m, n)}}$  and

- i. Every  $A_{(m, n)}$  is a sofic.
- ii.  $\lim_{n \rightarrow \infty} h(A_{(m, n)}) = h(\Gamma_\alpha)$ .

**Proof.** By the fact that for each  $n \in \mathbb{N}$ ,  $C_{u_1} \cup \dots \cup C_{u_n}$  is a finite graph, (i) is trivial.

For (ii), let  $\epsilon > 0$ . So there is  $n \geq 1$  such that

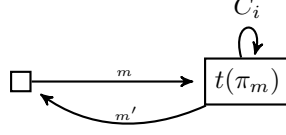
$$h(\Gamma_\alpha) - \epsilon < \frac{1}{n} \log |\{C : C \text{ is a cycle in } \Gamma_\alpha \text{ passing through } t(\pi_m), |C| = n\}|.$$

Let  $\{C : C \text{ is a cycle in } \Gamma_\alpha \text{ passing through } t(\pi_m), |C| = n\} = \{C_1, \dots, C_N\}$ . For each  $1 \leq i \leq N$ , set  $v_i := \mathcal{L}(C_i)m'$  Figure 6. So there is a cycle  $C'_i$  in  $\Gamma_\alpha$  labeled  $\mathcal{L}(C_i)m'm$  and so  $C'_i = C_{u_{j_i}}$  for some  $j_i \in \mathbb{N}$ . Thus  $C_1 \cup C_2 \cup \dots \cup C_N \subseteq \Gamma_{(m, t)}$  where  $t := \max\{j_1, j_2, \dots, j_N\}$ . Thus

$$N \leq |\{C : C \text{ is a cycle in } \Gamma_{(m, t)} \text{ passing through } t(\pi_m), |C| = n\}|$$

and so  $h(\Gamma_\alpha) - \epsilon < h(A_{(m, t)})$ . Hence  $h(\Gamma_\alpha) < \lim h(A_{(m, n)})$  and we are done.  $\square$  Let  $G$  be generating graph for a subshift  $X \subseteq \mathcal{A}^{\mathbb{Z}}$ . Fix

$\{a_1, a_2, \dots, a_r\} \subseteq \mathcal{A}$  and  $\{u_1, u_2, \dots, u_r\} \subseteq W(X)$ . We construct a new graph from  $G$  denoted by  $G_{u_i \leftrightarrow a_i}$  by replacing  $u_i$  for  $a_i$  whenever there is a path in  $G$  labeled  $a_i$  for all  $1 \leq i \leq r$ .



**Figure 6:** The graph  $G$ .

**Proposition 3.14.** Let  $\mathcal{A} = \{a_1, a_2, \dots, a_r\}$  and  $X_0^+$  be the Fischer cover of  $\mathcal{A}^{\mathbb{Z}}$  at the base point  $I_0$ . Pick  $\{u_1, u_2, \dots, u_r\} \subseteq W(X)$  such that for each  $1 \leq i \leq r$ ,

$$u_i \notin \{u : u \text{ is a finite concatenation of } \{u_j : 1 \leq j \leq r \text{ and } j \neq i\}\}.$$

For each  $N \in \mathbb{N}$ , set

$$\mathcal{N} := \{(n_1, n_2, \dots, n_r) \in (\mathbb{N} \cup \{0\})^r : k_1 n_1 + k_2 n_2 + \dots + k_r n_r = N\}$$

where  $k_i := |u_i|$ . Then,

$$h((X_0^+)_{u_i \leftrightarrow a_i}) = \lim_{N \rightarrow \infty} \frac{1}{N} \log \sum_{(n_1, n_2, \dots, n_r) \in \mathcal{N}} \binom{n}{n_1} \binom{n-n_1}{n_2} \dots \binom{n-n_1-\dots-n_{r-1}}{n_r}$$

where  $n = n_1 + n_2 + \dots + n_r$ .

**Proof.** For each  $1 \leq i \leq r$ , there is exactly one cycle labeled  $u_i$ . So

$$\sum_{(n_1, n_2, \dots, n_r) \in \mathcal{N}} \binom{n}{n_1} \binom{n-n_1}{n_2} \dots \binom{n-n_1-\dots-n_{r-1}}{n_r}$$

is the number of cycles of length  $N$  based at  $I_0$  and we are done.  $\square$

**Corollary 3.15.** Let  $m \in S_t(X)$ . For each  $u \in W(X_{(\alpha, 0)})$ , set  $k_i := |u_i u'_i m u''_i|$ . If all  $A_{(m, n)}$  is an SFT, then

$$h(\Gamma_\alpha) = \lim_{N \rightarrow \infty} \frac{1}{N} \log \sum_{(n_1, n_2, \dots, n_r) \in \mathcal{N}} \binom{n}{n_1} \binom{n-n_1}{n_2} \dots \binom{n-n_1-\dots-n_{r-1}}{n_r}.$$

For instance let  $X$  be the golden mean shift that formed of  $\{0, 1\}^{\mathbb{Z}}$  by replacing  $00$  instead  $0$ . Then,  $h(X) = \lim_{N \rightarrow \infty} \frac{1}{N} \log \sum_{n+2m=N} \binom{n+m}{n}$ .

## 4 Strong Synchronized Derived

Suppose  $\alpha \in S(X)/\sim$ . Let  $X_{(\alpha_s, 0)}$  denote the set of elements  $x \in X_{(\alpha, 0)}$  satisfying the following.

- i. If  $i \in \mathbb{Z}$ , then there are  $m, m' \in \alpha \cap S_t(X)$  such that  $m \subseteq x_{(-\infty, i]}$ ,  $m' \subseteq x_{[i, +\infty)}$ .
- ii. There is  $M > 0$  such that if  $u$  and  $v$  are the consecutive minimal pairs of  $\alpha \cap S_t(X)$  in  $x$ , then  $\text{gap}(u, v) \leq M$ .

$X_{(\alpha_s, 0)}$  is clearly a subshift of  $X_{(\alpha, 0)}$ , but generally not closed. Clearly  $\overline{X_{(\alpha_s, 0)}} = \overline{X_{(\alpha, 0)}}$  and

$$X_{(\alpha_s, 0)} \subseteq X_{(\alpha, 0)}. \quad (5)$$

The next example shows that the inclusion in (5) is not equality.

**Example 4.1.** In the golden mean shift  $X$ ,  $X = X_{(\alpha, 0)}$  and  $X_{(\alpha, 0)} - X_{(\alpha_s, 0)} = \{0^\infty\}$ .

If  $X$  is a synchronized (resp. strong synchronized) system, then there is exactly one  $\alpha \in S(X)/\sim$  such that  $\overline{X_{(\alpha, 0)}} = X$  (resp.  $\overline{X_{(\alpha_s, 0)}} = X$ ) ([9, Lemma 3.5]) and call it the *top component* (resp. *top strong component*).

Let  $X$  be a shift space and  $\alpha \in S(X)/\sim$ . Suppose  $(\overline{X_{(\alpha, 0)}})_c$  be the top component of the synchronized system  $\overline{X_{(\alpha, 0)}}$ . Thomsen in ([9]) prove that

$$X_{(\alpha, 0)} \subseteq (\overline{X_{(\alpha, 0)}})_c. \quad (6)$$

The next Proposition shows that for strong synchronized systems the inclusion in (6) is equality.

**Proposition 4.2.** *Let  $X$  be a shift space and  $\alpha \in S(X)/\sim$ . Suppose  $(\overline{X_{(\alpha_s, 0)}})_c$  be the strong top component of the strong synchronized system  $\overline{X_{(\alpha_s, 0)}}$ . Then,  $X_{(\alpha_s, 0)} = (\overline{X_{(\alpha_s, 0)}})_c$ .*

**Proof.** By definition 3.2,  $S_t(\overline{X_{(\alpha, 0)}}) \subseteq S_t(X)$  and so  $X_{(\alpha_s, 0)} = (\overline{X_{(\alpha_s, 0)}})_c$ .  $\square$

Let  $X$  be an strong synchronized system. We set

$$\partial_s X = \{x \in X \mid u \subseteq x \Rightarrow u \notin S_t(X)\}.$$

and call it the *strong derived* shift space of  $X$ . Since  $\partial_s X$  is a shift space we can continue, and consider  $\partial_s(\partial_s X) = \partial_s^2 X$ ,  $\partial_s(\partial_s^2 X) = \partial_s^3 X$ , etc. Of course, it can happen that these constructions give nothing interesting; it can be that there are no synchronizing blocks for  $X$ , in which case  $\partial_s X = X$ . For convenience we set  $\partial_s^0 X = X$ . We define the *strong depth* of  $X$  to be

$$\text{Depth}_s(X) = \sup\{n \in \mathbb{N} : \partial_s^n X \neq \emptyset\}.$$

Thus a minimal shift space with infinitely many points as well as an SFT have depth 0, but for different reasons.

**Example 4.3.** i. In example 3.3,  $\partial_s S = S$  and  $\partial S = X$ .

ii. Let  $Y$  be a sofic. Then, by 3.4  $\partial_s Y = \partial Y$  and so by ([9, Theorem 6.6]),  $\partial_s Y$  is sofic.

iii. Suppose  $Z$  be synchronized system. Then, by ([8, Theorem 6.16]),  $h(Z) = \max\{h(X_0^+), h(\partial_s Z)\}$ .

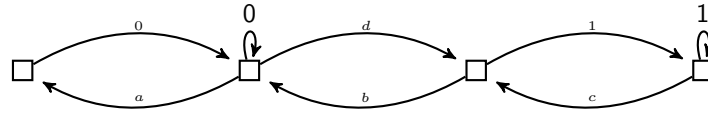
**Proposition 4.4.** *If  $m \in S(X)$  and  $m \notin S_t(X)$ , then  $m$  is a synchronizing block of  $\partial_s X$ .*

**Proof.** Let  $am, mb \in W(\partial_s X)$ . So there are  $x, y \in \partial_s X$  such that  $am = x_{[i, 0]}$  and  $mb = y_{(0, j]}$  for some  $i, j \in \mathbb{Z}$ . Since  $m \in S(X)$ , so  $x_{(-\infty, i)} amby_{(j, +\infty)} \in X$ . It suffice to show that

$$S_t(X) \cap \{x_{[i_0, j_0]} : i_0 \leq 0, j_0 > 0\} = \emptyset.$$

Pick  $i_0 \leq 0$  and  $j_0 > 0$ . There is unique vertex  $I$  of Fischer cover  $\Gamma_{[m]}$  with this properties  $m \in F_-(I)$  and  $w_+(I) = w_+(m)$ . Since  $x_{[i_0, 0]} \notin S_t(X)$ , so there are two finite paths in  $\Gamma_{[m]}$  labeled  $x_{[i_0, 0]}$  and terminating at  $I$  and so  $x_{[i_0, j_0]} \notin S_t(X)$ .  $\square$  The next example shows that there are systems  $X$  such that  $m \in S(X) - S_t(X)$  which may not be a synchronized.

**Example 4.5.** Let  $G$  be the graph in Figure 7 and  $X = X(G)$ . Then,  $\partial_s X = \{0^\infty, 1^\infty\}$ .



**Figure 7:** The graph  $G$ .

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